

The aim of this paper is to describe the distribution of definable elements in models of PA. The analysis of definable elements proved to be fruitful in the context of investigating the relations between various types of induction, collection and minimalization. The classical results in this area were obtained by Kirby and Paris in (5); see also (4) for the discussion of parameter free induction. In addition, (2) and (3) can be recommended as a very useful, general reference. For more results in this direction, see (1).

We start with the following definition of structures consisting of elements of a model  $M$  definable by formulas belonging to  $\Sigma_n$  ( $\Pi_n$ ).<sup>1</sup>

**Definition 1.**

$$K^{\Sigma_n}(M) = \{a \in M : \exists \varphi(x) \in \Sigma_n \varphi(x) \text{ defines } a \text{ in } M\}$$

$$K^{\Pi_n}(M) = \{a \in M : \exists \varphi(x) \in \Pi_n \varphi(x) \text{ defines } a \text{ in } M\}$$

The first useful fact concerns the relation between  $Th_{\Pi_n}(N)$  (the set of all  $\Pi_n$  sentences true in the standard model) and  $Th_{\Sigma_{n+1}}(N)$  (an analogous set of  $\Sigma_{n+1}$  sentences).

**Fact 1.**  $\forall n \in \mathbb{N} \ Th_{\Pi_n}(N) \vdash Th_{\Sigma_{n+1}}(N)$

*Proof.* Let  $M \models Th_{\Pi_n}(N)$ . Let  $\varphi$  be a  $\Sigma_{n+1}$  formula true in  $N$  of the form  $\ulcorner \exists x \psi(x) \urcorner$ . So for some  $n$  belonging to  $N$ ,  $N \models \psi(n)$ . Since  $\psi(n)$  belongs to  $\Pi_n$ , we have:  $M \models \psi(n)$ ; therefore  $M \models \varphi$ .

□

The next two theorems state the sufficient and necessary conditions of the existence of  $\Sigma_n$  (or  $\Pi_n$ ) definable, nonstandard elements in models of PA. In what follows, by „ $K^{\Sigma_n}(M) > N$ ” we mean „there are nonstandard,  $\Sigma_n$ -definable elements of  $M$ ”; analogously for  $K^{\Pi_n}(M)$ .<sup>2</sup>

**Theorem 1.** For every  $n > 0$ , for every  $M \models \text{PA}$ ,  $K^{\Sigma_n}(M) > N$  iff  $M \not\models Th_{\Pi_n}(N)$ .

*Proof.* Assume at first, that  $\varphi(x)$  is a  $\Sigma_n$  formula which defines in  $M$  a nonstandard element. Assume also that  $M \models Th_{\Pi_n}(N)$ . Then we have:  $N \models \exists x \varphi(x)$ , because otherwise  $N \models \forall x \neg \varphi(x)$ , so we would obtain:  $M \models \forall x \neg \varphi(x)$ , since this last formula belongs to  $\Pi_n$  (assuming that  $n > 0$ ). So we obtain a certain  $k \in \mathbb{N}$  such that  $N \models \varphi(k)$ , but  $M \not\models \varphi(k)$ , since  $\varphi(x)$  defines something nonstandard in  $M$ . However,  $\varphi(k) \in \Sigma_n$ , therefore  $M \models \varphi(k)$  by our

<sup>1</sup> For the definition of the classes  $\Sigma_n$  and  $\Pi_n$ , see e.g. (2) or (3).

<sup>2</sup> Theorem 1 is folklore; we owe the idea of the proof of theorem 2 to H. Kotlarski and K. Zdanowski.

assumption that  $M \not\models Th_{\Pi_n}(N)$  (cf. Fact 1). This is impossible, since  $\varphi(x)$  was to define something nonstandard in  $M$ .

For the second implication, fix the smallest  $k$  such that  $M \in Th_{\Pi_n}(N)$  (then  $k \leq n$ ). Let  $\psi = \ulcorner \forall x \varphi(x) \urcorner$  with  $\varphi(x)$  belonging to  $\Sigma_{k-1}$ , such that  $N \models \forall x \varphi(x)$  and  $M \models \exists x \neg \varphi(x)$ . Let  $\tau(x)$  be a  $\Sigma_k$  formula which states that  $x$  is the smallest number satisfying  $\neg \varphi(x)$ . Then  $\tau(x)$  defines something nonstandard in  $M$ , because if  $M \models \tau(n)$  for some  $n \in N$ , then  $M \models \neg \varphi(n)$ , but  $N \models \varphi(n)$ , so  $\varphi(n) \in Th_{\Sigma_{k-1}}(N)$ . Therefore  $M \not\models Th_{\Pi_{k-1}}(N)$ , because otherwise  $M \models \varphi(n)$ . We obtain a contradiction, because  $k$  was to be the smaller number of this kind. □

**Theorem 2.** For every  $n \geq 0$ , for every  $M \models \text{PA}$ ,  $K^{\Pi_n}(M) > N$  iff  $M \not\models Th_{\Pi_{n+1}}(N)$ .

*Proof.*

( $\rightarrow$ ) Assume that  $\varphi(x)$  is a  $\Pi_n$  formula which defines in  $M$  a nonstandard element. Assume also that  $M \models Th_{\Pi_{n+1}}(N)$ . Then we have:  $N \models \exists x \varphi(x)$ , because otherwise  $N \models \forall x \neg \varphi(x)$ . Since this last formula is  $\Pi_{n+1}$ , we would have:  $M \models \forall x \neg \varphi(x)$ , which is impossible. So let  $k \in N$  be such that  $N \models \varphi(k)$ . Therefore  $M \models \varphi(k)$ ; but this is a contradiction, since  $\varphi(x)$  was to define something nonstandard in  $M$ .

( $\leftarrow$ ) Let  $n$  be the smallest natural number such that  $M \not\models Th_{\Pi_{n+1}}(N)$ . If  $n = 0$ , we have:  $N \models \forall x \varphi(x)$  and  $M \models \exists x \neg \varphi(x)$  for some formula  $\varphi(x)$  belonging to  $\Delta_0$ . Then the smallest  $x$  such that  $\neg \varphi(x)$  is a nonstandard,  $\Delta_0$  definable element of  $M$ .

So now let's assume that  $n > 0$ . Let  $N \models \forall x \exists y \neg \varphi(y, x)$  and  $M \models \exists x \forall y \varphi(y, x)$  for  $\varphi$  belonging to  $\Sigma_{n-1}$ . In this case for every  $a$ , if  $M \models \forall y \varphi(y, a)$ , then  $a > N$ , because otherwise we would have for  $k \in N$ :  $M \models \forall y \varphi(y, k)$ . But  $N \models \exists y \neg \varphi(y, k)$  and since this last formula is  $\Sigma_n$  it turns out that  $M \not\models Th_{\Pi_n}(N)$ , therefore  $M \not\models Th_{\Pi_n}(N)$ , which contradicts the choice of  $n$ .

Now take the following formula  $\tau(z)$  belonging to  $\Pi_n$ :

$$\begin{aligned} \forall a, b [z = (a, b) \Rightarrow & (\forall y \varphi(y, a) \wedge \\ & \wedge \forall w < a \exists y < b \neg \varphi(y, w) \wedge \\ & \wedge \forall c < b \exists w < a \forall y < c \varphi(y, w))] \end{aligned}$$

So  $\tau(z)$  states in effect, that  $z$  is a pair with the following properties:

- (i)  $z$  is a pair whose first element  $a$  is a witness for  $\ulcorner \exists x \forall y \varphi(y, x) \urcorner$
- (ii) the second element  $b$  of  $z$  is an upper bound for counterexamples below  $a$
- (iii)  $b$  is the smallest such upper bound.

The first two conditions uniquely identify  $a$ , the last one uniquely identifies  $b$ , so in effect  $\tau(z)$  is a  $\Pi_n$  formula defining a nonstandard element of  $M$ . □

Our next result is a solution to the following problem: in which models of PA are there elements which are essentially  $\Pi_n$ -definable; we mean by that the elements which are  $\Pi_n$ , but not  $\Sigma_n$  definable.<sup>3</sup> In effect, our result is a partial answer to the question if the  $\Sigma_n$ - $\Pi_n$  hierarchy is strict.

**Theorem 3.** For every  $n$ , for every  $M \models \text{PA}$

If  $M \neq \text{Th}_{\Pi_{n+1}}(N)$ , then for every  $k \geq n+1$   $K^{\Sigma_k}(M) \subset K^{\Pi_k}(M)$ .

*Proof.* Fix  $M$ ,  $n$  and  $k \geq n+1$ . Let  $\psi(x, y)$  be the following formula belonging to  $\Pi_k$ :

$$\begin{aligned} \forall \varphi \leq x [(\varphi \in \Pi_{k-1} \wedge \exists w \text{Tr}_{\Pi_{k-1}}(\varphi(w))) \Rightarrow \exists w < y \text{Tr}_{\Pi_{k-1}}(\varphi(w))] \wedge \\ \wedge \exists \varphi \leq x [\exists w \leq y \text{Tr}_{\Pi_{k-1}}(\varphi(w)) \wedge \forall w < y-1 \neg \text{Tr}_{\Pi_{k-1}}(\varphi(w))] \end{aligned}$$

where  $\text{Tr}_{\Pi_{k-1}}(x)$  is a universal formula for  $\Pi_{k-1}$  sentences. The formula  $\psi(x, y)$  states in effect, that  $y$  is the least upper bound for witnesses for  $\Pi_{k-1}$  formulas below  $x$ . By Theorem 1, let  $\gamma(x)$  be a  $\Sigma_k$  definition of some nonstandard element of  $M$  (here we use the assumption that  $M \neq \text{Th}_{\Pi_{n+1}}(N)$ ). Then the formula

$$\forall x [\gamma(x) \Rightarrow \psi(x, y)]$$

is a  $\Pi_k$  definition of some element  $a$ . Now assume that  $a$  is  $\Sigma_k$  definable by a formula  $\ulcorner \exists z \chi(z, \cdot) \urcorner$ . We take the following formula  $\xi(w)$  belonging to  $\Pi_{k-1}$ :

$$\exists w_1, w_2 < w [w = (w_1, w_2) \wedge \chi(w_1, w_2)]$$

We have:  $M \models \exists w \xi(w)$ , so by our choice of  $a$ ,  $M \models \exists w < a \xi(w)$ . But any such  $w$  would be a pair whose second element is  $a$ , since  $\ulcorner \exists z \chi(z, \cdot) \urcorner$  defines  $a$ . This is impossible and we obtain a contradiction.

□

Note that the converse to Theorem 3 is false. As a counterexample, we may take  $M \models \text{Th}_{\Pi_{n+1}}(N)$  such that  $M \neq \text{Th}_{\Pi_{n+2}}(N)$ . Then by Theorem 2,  $K^{\Pi_{n+1}}(M) > N$ , but by Theorem 1,  $K^{\Sigma_{n+1}}(M) = N$ , so  $K^{\Sigma_{n+1}}(M) \subset K^{\Pi_{n+1}}(M)$ . By Theorem 3, we have also  $K^{\Sigma_k}(M) \subset K^{\Pi_k}(M)$  for all  $k \geq n+2$ . So the subsequent of Theorem 3 holds, but the antecedent is obviously false.

In a similar vein, we consider now the question: are there  $\Sigma_{n+1}$  definable elements which are not  $\Pi_n$  definable. Together with the previous result our next theorem gives us in effect the strict  $\Sigma_n$ - $\Pi_n$  hierarchy. However, Theorem 4 contains also some extra information: there is always an essentially  $\Sigma_{n+1}$  definable element below an arbitrarily chosen  $\Sigma_{n+1}$  definable one.

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<sup>3</sup> Note that a  $\Sigma_n$ -definable element is always  $\Pi_n$ -definable, because if  $\varphi(x)$  is  $\Sigma_n$  and defines  $a$ , then  $\ulcorner \forall x [\varphi(x) \Rightarrow y = x] \urcorner$  is  $\Pi_n$  and defines  $a$ .

**Theorem 4.** Let  $M \models \text{PA}$ , let  $n \in \mathbb{N}$  and let  $a \in M$  such that  $a > N$  and  $a \in K^{\Sigma_{n+1}}(M)$ . Then for some  $b < a$ ,  $b \in K^{\Sigma_{n+1}}(M)$  and  $b \notin K^{\Pi_n}(M)$ .

*Proof.* Fix an appropriate  $M$ ,  $n$  and  $a$ . Let  $\psi(z) \in \Sigma_{n+1}$  and define  $a$ . Consider the following formula  $\chi(s)$  belonging to  $\Sigma_{n+1}$ :

$$\begin{aligned} \exists z [\psi(z) \wedge \forall \varphi < z (\varphi \in \Pi_n \Rightarrow (\neg \text{Tr}_{\Pi_n}(\varphi(s)) \vee \exists w < s \text{Tr}_{\Pi_n}(\varphi(w)))) \wedge \\ \wedge \forall w < s \exists \varphi < z [\varphi \in \Pi_n \wedge \text{Tr}_{\Pi_n}(\varphi(w)) \wedge \forall x < w \neg \text{Tr}_{\Pi_n}(\varphi(x))] \end{aligned}$$

Intuitively, the formula  $\chi(s)$  states that:

1.  $s$  is not the smallest object satisfying some  $\Pi_n$  formula below  $a$
2.  $s$  is the smallest object with the above property.

Now we are going to show that  $\chi(s)$  defines something below  $a$ . Let  $\alpha(\varphi, s)$  be the formula:

$$\neg \text{Tr}_{\Pi_n}(\varphi(s)) \vee \exists w < s \text{Tr}_{\Pi_n}(\varphi(w))$$

So the formula  $\alpha(\varphi, s)$  states that  $\varphi$  doesn't identify  $s$ .

In order to see that  $\chi(s)$  defines something below  $a$ , it is clearly enough to observe that:

$$\text{PA} \vdash \forall z \exists s \leq z \forall \varphi < z [\varphi \in \Pi_n \Rightarrow \alpha(\varphi, s)]$$

For an indirect proof, fix  $z$  in a model  $M$  of PA such that in  $M$ :

$$\forall s \leq z \exists \varphi < z [\varphi \in \Pi_n \wedge \neg \alpha(\varphi, s)]$$

From pigeonhole principle we obtain:

$$\exists s_1, s_2 \leq z \exists \varphi < z [s_1 \neq s_2 \wedge \varphi \in \Pi_n \wedge \neg \alpha(\varphi, s_1) \wedge \neg \alpha(\varphi, s_2)]$$

But this means that both  $s_1$  and  $s_2$  are the smallest numbers satisfying  $\varphi$ , but nonetheless they are different. Since this is impossible, it ends the proof.

□

Now we are ready to consider in general terms the question concerning the distribution of definable elements in models of PA. In accordance with the usual convention, we write " $\subseteq_e$ " (" $\subseteq_{cf}$ ") to denote an end (cofinal) extension of a given structure.

**Theorem 5.** For every  $n$ , for every  $M \models \text{PA}$ :

- (a)  $K^{\Sigma_n}(M) \subseteq_e K^{\Pi_n}(M)$
- (b)  $K^{\Pi_n}(M) \subseteq_{cf} K^{\Sigma_{n+1}}(M)$

*Proof.*

- (a) Let  $\varphi(x) \in \Sigma_n$  and define  $a$ . Then the formula:  $\lceil \forall x[\varphi(x) \Rightarrow x = y] \rceil$  is  $\Pi_n$  and defines  $a$ . This shows that  $K^{\Sigma_n}(M) \subseteq K^{\Pi_n}(M)$ . As for being an initial segment, let  $\varphi(x) \in \Pi_n$  and define  $a$ ; let  $\psi(x) \in \Sigma_n$  and define  $b$ . Assume also that  $a < b$ . Then  $a$  is  $\Sigma_n$  definable by the formula:

$$\exists x [\psi(x) \wedge \forall w < x(\varphi(w) \Rightarrow w = y)]$$

- (b) The inclusion itself is obvious. So we show only that there is a  $\Pi_n$  definable element above each  $\Sigma_{n+1}$  definable one.

Let  $a \in K^{\Sigma_{n+1}}(M)$ . Let  $\psi(x)$  be a  $\Sigma_{n+1}$  definition of  $a$ . If  $n = 0$ , then  $\psi(x)$  has a form  $\lceil \exists y\varphi(y, x) \rceil$  for  $\varphi$  belonging to  $\Delta_0$ . In this case the formula “ $w$  is the smallest pair such that  $\varphi(w_0, w_1)$ ” is  $\Delta_0$  and defines a number greater than  $a$ .

So now we may assume that  $n > 0$  and our formula  $\psi(x)$  has the form:  $\lceil \exists s\forall y\varphi(y, s, x) \rceil$ . Then  $M \models \exists x[x = (x_0, x_1) \wedge \forall y\varphi(y, x_0, x_1)]$ . Now let  $\gamma(z)$  be the following formula:

$$\begin{aligned} \exists c, d < z [z = (c, d) \wedge \forall y\varphi(y, c_0, c_1) \wedge \\ \forall w < c \exists y < d \neg \varphi(y, w_0, w_1) \wedge \\ \forall e < d \exists w < c \forall y < e \varphi(y, w_0, w_1)] \end{aligned}$$

(In our notation e.g. “ $w_0$ ” denotes the left element of the pair  $w$ ).

This formula defines a number  $z$ , which is a pair  $(c, d)$ , where  $c = (\cdot, a)$ , therefore  $z > a$ . In addition,  $z \in K^{\Pi_n}(M)$ , because  $\gamma(z)$  is a  $\Pi_n$  formula.

□

The following corollary states the sufficient condition under which the above inclusions are proper.

**Corollary 1.** Let  $M \not\equiv Th_{\Pi_{n+1}}(N)$ . Then:

- (a)  $K^{\Sigma_{n+1}}(M) \subset_e K^{\Pi_{n+1}}(M)$   
(b)  $K^{\Pi_n}(M) \subset_{cf} K^{\Sigma_{n+1}}(M)$

**Proof.**

- (a) By Theorem 5a,  $K^{\Sigma_{n+1}}(M) \subset_e K^{\Pi_{n+1}}(M)$ . Since  $M \not\equiv Th_{\Pi_{n+1}}(N)$ , by Theorem 3 we obtain:  $K^{\Sigma_{n+1}}(M) \subset_e K^{\Pi_{n+1}}(M)$ , which is our desired result.  
(b) We have:  $K^{\Pi_n}(M) \subset_{cf} K^{\Sigma_{n+1}}(M)$  by Theorem 5b. Since  $M \not\equiv Th_{\Pi_{n+1}}(N)$ , we obtain  $K^{\Sigma_{n+1}}(M) > N$  by Theorem 1. So by Theorem 4,  $K^{\Pi_n}(M) \subset_{cf} K^{\Sigma_{n+1}}(M)$ . And that ends the proof.

□

Corollary 2 formulated below is a negative result. It states that no interesting relations (no relations of being an initial segment or having a cofinal extension) hold between  $\Pi_n$ -type structures, unless in the standard case.

**Corollary 2.**

- (a)  $K^{\Pi_n}(M) \subseteq_e K^{\Pi_{n+1}}(M)$  iff  $K^{\Pi_n}(M) = N$
- (b)  $K^{\Pi_n}(M) \subseteq_{cf} K^{\Pi_{n+1}}(M)$  iff  $K^{\Pi_{n+1}}(M) = N$

*Proof.*

- (a) The implication from right to left is obvious. So let's assume that  $K^{\Pi_n}(M) \neq N$ . In this case by Theorem 2,  $M \not\equiv Th_{\Pi_{n+1}}(N)$ . So by Corollary 1b, there is an  $a \in K^{\Sigma_{n+1}}(M)$  such that  $a \notin K^{\Pi_n}(M)$ . Therefore  $a \in K^{\Pi_{n+1}}(M)$ . But  $K^{\Pi_n}(M) \subseteq_{cf} K^{\Sigma_{n+1}}(M)$  (Theorem 5b), so there is a  $b$  belonging to  $K^{\Pi_n}(M)$  such that  $b > a$ . Therefore it is not the case that  $K^{\Pi_n}(M) \subseteq_e K^{\Pi_{n+1}}(M)$ .
- (b) The implication from right to left is obvious. So let's assume that  $K^{\Pi_n}(M) \subseteq_{cf} K^{\Pi_{n+1}}(M)$  and that  $K^{\Pi_{n+1}}(M) \neq N$ . So  $K^{\Pi_n}(M) \neq N$ . By Theorem 2,  $M \not\equiv Th_{\Pi_{n+1}}(N)$ . Therefore by Corollary 1a,  $K^{\Sigma_{n+1}}(M) \subset_e K^{\Pi_{n+1}}(M)$ , so we obtain: not  $K^{\Pi_n}(M) \subseteq_{cf} K^{\Pi_{n+1}}(M)$ , which ends the proof, producing the desired contradiction.

□

As we already remarked, structures on  $\Sigma_n$ -definable elements proved to be useful in analysing the relations between various types of induction. A typical question worth considering in this context was: how much induction satisfies a given  $\Sigma_n$ -structure. Obviously one can ask the same question about  $\Pi_n$ -structures. However, the next lemma leads to a negative conclusion: such structures don't even satisfy  $PA^-$  (arithmetic without induction scheme). On a more positive vein, the lemma will permit us to characterize  $K^{\Sigma_{n+1}}(M)$  as the closure of  $K^{\Pi_n}(M)$  under subtraction (which is Corollary 3).

**Lemma 1.** For every  $n \geq 0$ , for every  $w \in K^{\Sigma_{n+1}}(M)$ , there is an  $a \in K^{\Pi_n}(M)$  such that  $w+a \in K^{\Pi_n}(M)$ .

*Proof.* Fix  $w \in K^{\Sigma_{n+1}}(M)$ . Let  $\ulcorner \exists s \forall y \varphi(y, s, x) \urcorner$  be a  $\Sigma_{n+1}$  definition of  $w$ . Pick an object  $a$  as in the proof of Theorem 5b - then  $a$  is a pair  $(c, d)$  such that  $(c_0, c_1)$  satisfy  $\ulcorner \forall y \varphi(y, s, x) \urcorner$  and  $d$  is the smallest upper bound for counterexamples below  $c$ . Such an  $a$  belongs to  $K^{\Pi_n}(M)$ . Then  $w+a$  has the following  $\Pi_n$  definition:

$$\exists xy < s[x = a \wedge \exists z < x(z = a_0 \wedge y = z_1 \wedge s = x + y)]$$

Remember that  $\ulcorner x = a \urcorner$  can be written down as a  $\Pi_n$  formula without parameters. As usual, we write  $a_0$  and  $a_1$  to denote the left and the right member of the pair  $a$ .

□

**Corollary 3.**  $K^{\Sigma_{n+1}}(M)$  is exactly the closure of  $K^{\Pi_n}(M)$  under subtraction.

*Proof.* Obvious, from Lemma 1.

□

It follows in particular that for every  $n \geq 0$ ,  $K^{\Pi_n}(M) \neq \text{PA}^-$ . Take simply an object  $w$  belonging to  $K^{\Sigma_{n+1}}(M)$  such that  $w \notin K^{\Pi_n}(M)$ . By Lemma 1, fix an  $a$  belonging to  $K^{\Pi_n}(M)$  such that  $w+a \in K^{\Pi_n}(M)$ . Therefore in  $K^{\Pi_n}(M)$   $w+a$  is not a sum of  $a$  with any other number.

Our next result is an addition to Theorem 4. We showed before that there is an essentially  $\Sigma_{n+1}$  definable element below each  $\Sigma_n$  definable one. Now we consider a special case of a  $\Pi_n$  definable element  $a$  greater than  $K^{\Sigma_n}(M)$ . We show that such an  $a$  determines an essentially  $\Sigma_{n+1}$ -definable  $b$  which similarly will be greater than  $K^{\Sigma_n}(M)$ .

In the proof we will use the following fact (we write “ $I^n(M)$ ” to denote an initial segment of  $M$  cofinal with  $K^{\Sigma_n}(M)$ ).<sup>4</sup>

**Fact 2.** For every  $n \geq 1$ ,  $I^n(M) \models \text{Th}_{\Pi_{n+1}}(M)$

For the proof, see (3).<sup>5</sup>

**Theorem 6.** Let  $n \geq 1$ . For every  $a \in K^{\Pi_n}(M)$ , if  $a > K^{\Sigma_n}(M)$ , then there is a number  $b < a$  such that  $b \in K^{\Sigma_{n+1}}(M)$  and  $b > K^{\Sigma_n}(M)$  and  $b \notin K^{\Pi_n}(M)$ .

*Proof.* Fix  $a \in K^{\Pi_n}(M)$  such that  $a > K^{\Sigma_n}(M)$ . Let  $s$  be the smallest number such that no  $\Pi_n$  formula below  $a$  identifies  $s$  (in other words, let  $s$  satisfy the formula  $\chi(s)$  used in the proof of Theorem 4). So  $s < a$ ,  $s \in K^{\Sigma_{n+1}}(M)$  and  $s \notin K^{\Sigma_n}(M)$ . We now claim that  $s > K^{\Sigma_n}(M)$ . By Fact 2,  $I^n(M) \models \text{Th}_{\Pi_{n+1}}(M)$ . Therefore for  $n \geq 1$ :

$$I^n(M) \models \forall x \exists \varphi [\varphi \in \Delta_0 \wedge \text{Tr}_0(\ulcorner \varphi(x) \urcorner) \wedge \forall w < x \neg \text{Tr}_0(\ulcorner \varphi(w) \urcorner)]$$

It follows that  $s > I^n(M)$ , therefore  $s > K^{\Sigma_n}(M)$ . Otherwise we would have:

$$M \models \exists \varphi < a [\varphi \in \Delta_0 \wedge s \text{ is the smallest object satisfying } \varphi]$$

But this is impossible by the choice of  $s$ .

□

## References

<sup>4</sup> Formally:  $I^n(M) = \{a \in M : \exists b \in K^{\Sigma_n}(M) M \models a \leq b\}$ .

<sup>5</sup> This is Theorem 10.8 in Kaye’s book.

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